

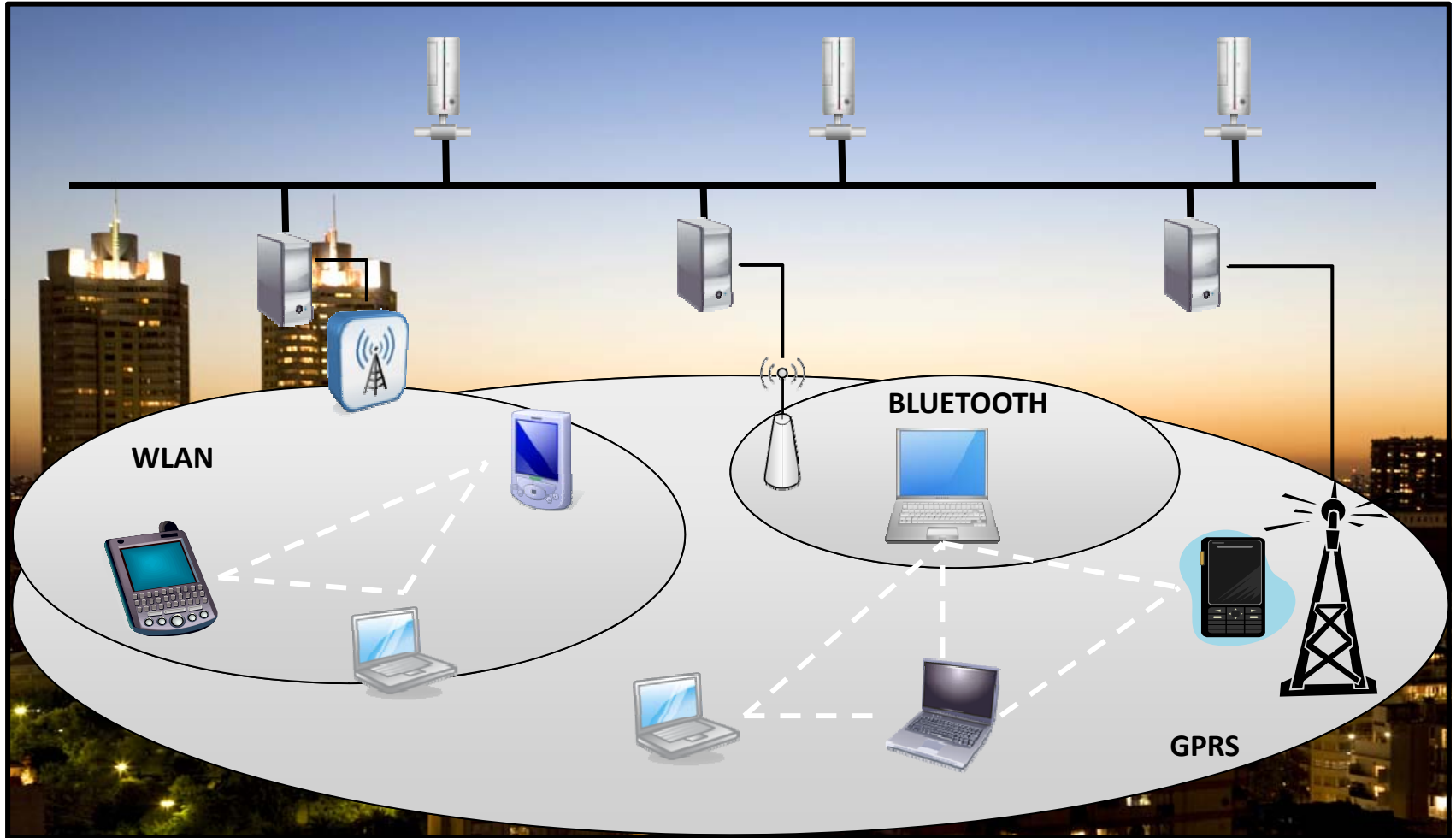
# Adaptive Context Data Distribution with Guaranteed Quality for Mobile Environments

Antonio Corradi, Mario Fanelli, Luca Foschini

*Dipartimento di Elettronica Informatica e Sistemistica -  
Università di Bologna*

- Context data distribution in mixed infrastructured/ad-Hoc environments
- Quality of Context (QoC)-based context data distribution
- Scalable context-Aware middleware for mobile Environments (SALES)
  - Data distribution process
  - Mapping QoC objectives to query parameters
- Implementation insights
- Experimental evaluations
- Lessons learned and ongoing work

# Context-aware applications in mobile and densely populated environments



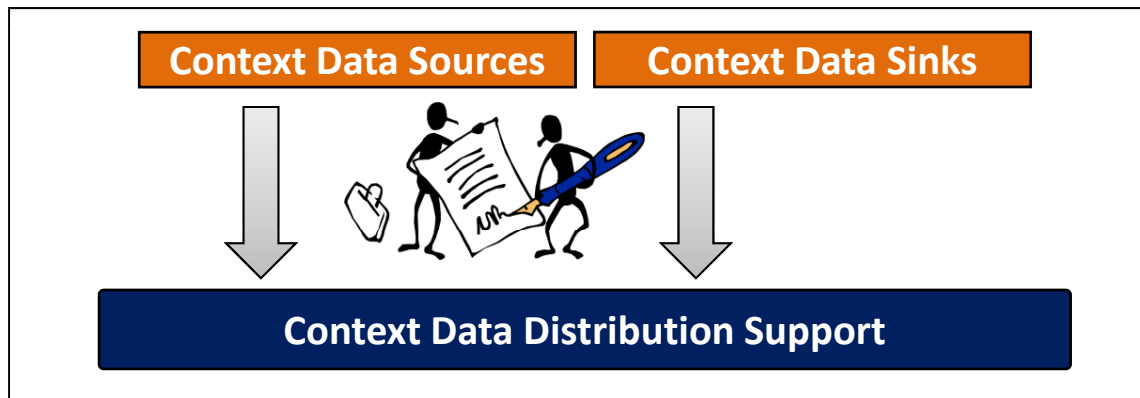
**Context data distribution** is a complex task that poses several challenging requirements:

- **Heterogeneity** of the computing environment: devices (smartphones, Personal Digital Assistants, netbooks, ...) and communication technologies (WiFi, Bluetooth, cellular 3G) and types (infrastructure and ad-hoc)
- **Device mobility** and **density**: ever-increasing number of mobile devices, already producing huge amounts of context data (environmental sensing, social computing, ...)
- **Data delivery with guaranteed quality levels**: depending on specific service, (disaster recovery and emergency scenario, entertainment, ...)

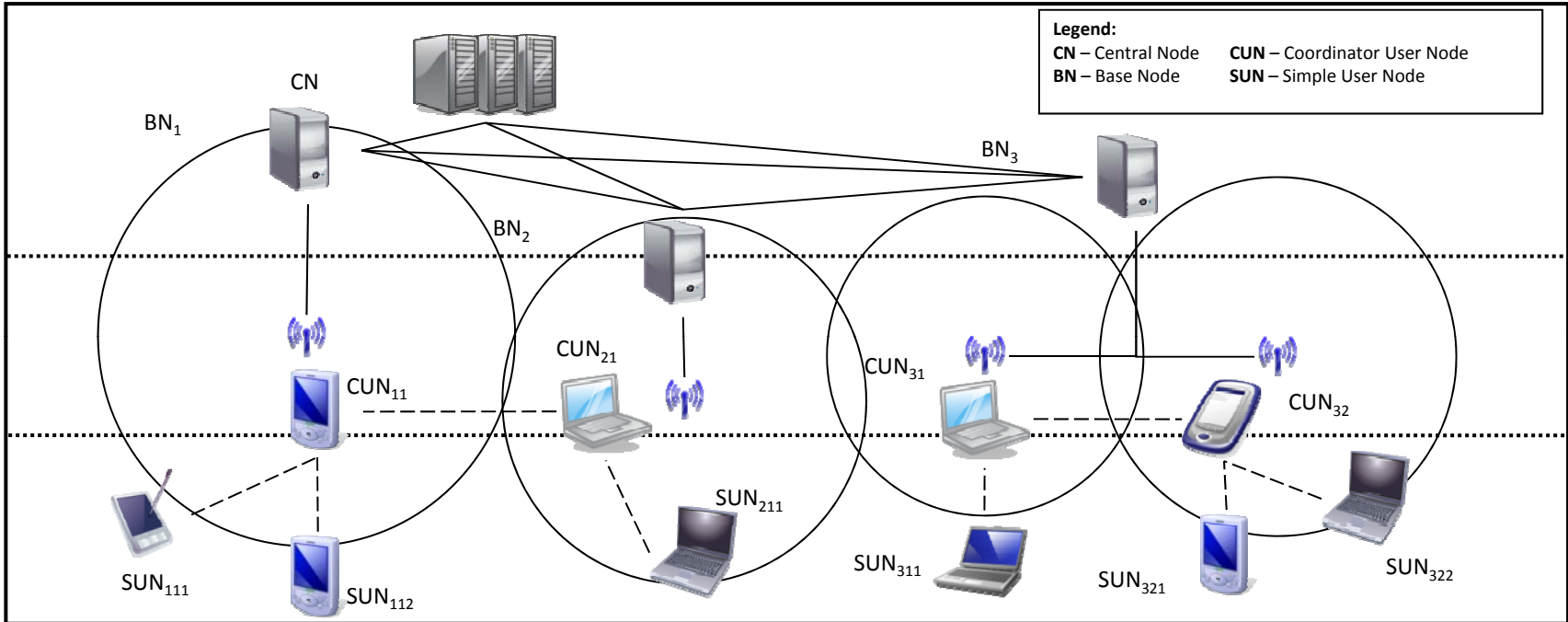
We claim the need for novel **Context Data Distribution Infrastructures (CDDIs)** to transparently address and take over context data distribution aspects (integration aspects, ***data distribution differentiation, scalability, ...***)

## Quality of Context (QoC)

- A concept usually applied to the quality of the distributed context data
- We extend this concept to tailor context data distribution and delivery so to save previous system-level resources

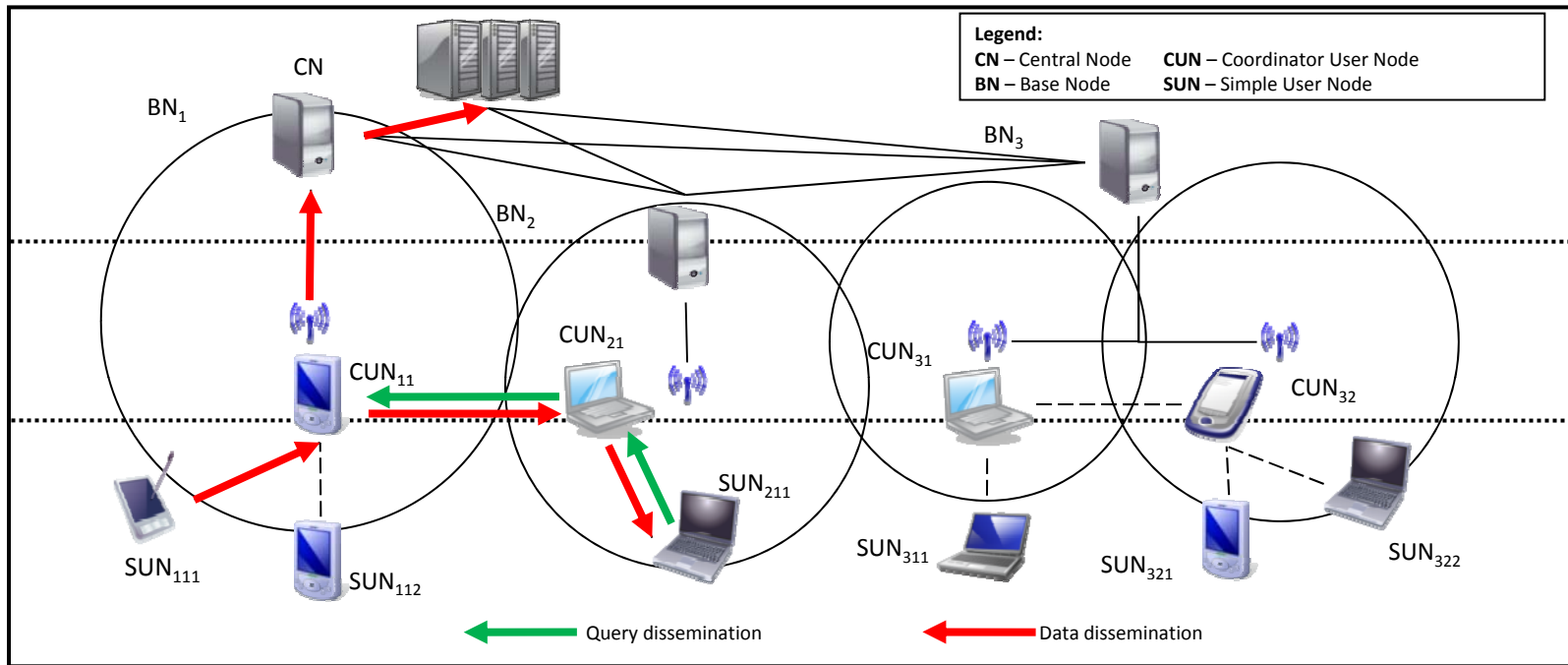


**Context Data Distribution Level Agreements (CDDLAs)**  
to detail the quality level the infrastructure has to ensure



## Three-level tree-like architecture

- minimizes tree depth to reduce management overhead
- ensures effective and integrated usage of wireless infrastructure and ad-hoc communication modes



- SALES distributes **context queries** to build dissemination paths
- Queries are disseminated both on the same level (horizontal distribution) and to the level above (vertical distribution)
- Data flow only on the bottom-up path between the data creator node and the CN
- Different dissemination paths are considered only when matching queries exist

SALES CDDLA comprises three main objectives:

1. **Freshness:** up-to-dateness requirement on matching data
2. **Data retrieval time:** the time needed by the mobile node to retrieve context data
3. **Priority:** traditional priority value used to enable traffic differentiation and to favor the routing of high-priority data under load conditions

To tailor query distribution, each SALES query has four main parameters:

1. **Horizontal time-to-live (HTTL):** used to limit the number of nodes traversed at the same hierarchy level
2. **Routing delay (RD):** used to delay query distribution to the next hop
3. **Query total lifetime (QTL):** used to handle query aging. When zero, the query is expired, and discarded by the system
4. **Query priority (QP):** used to enable priority-based query/data forwarding

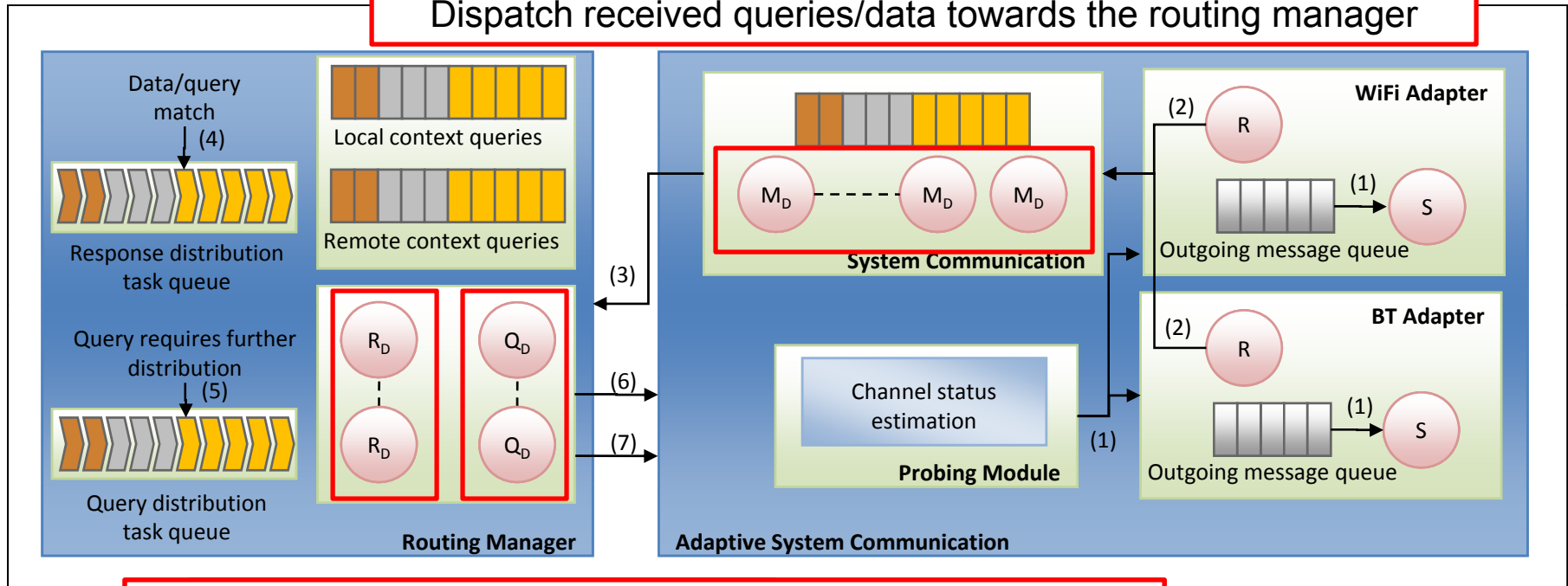


For the sake of clarity, we define three principal users' classes:

- **Gold** [latest version of data, retrieval time 2s, priority 0]
- **Silver** [valid data, retrieval time 4s, priority 1]
- **Bronze** [non-valid data expired at most by 2 seconds, retrieval time 6s, priority 2]

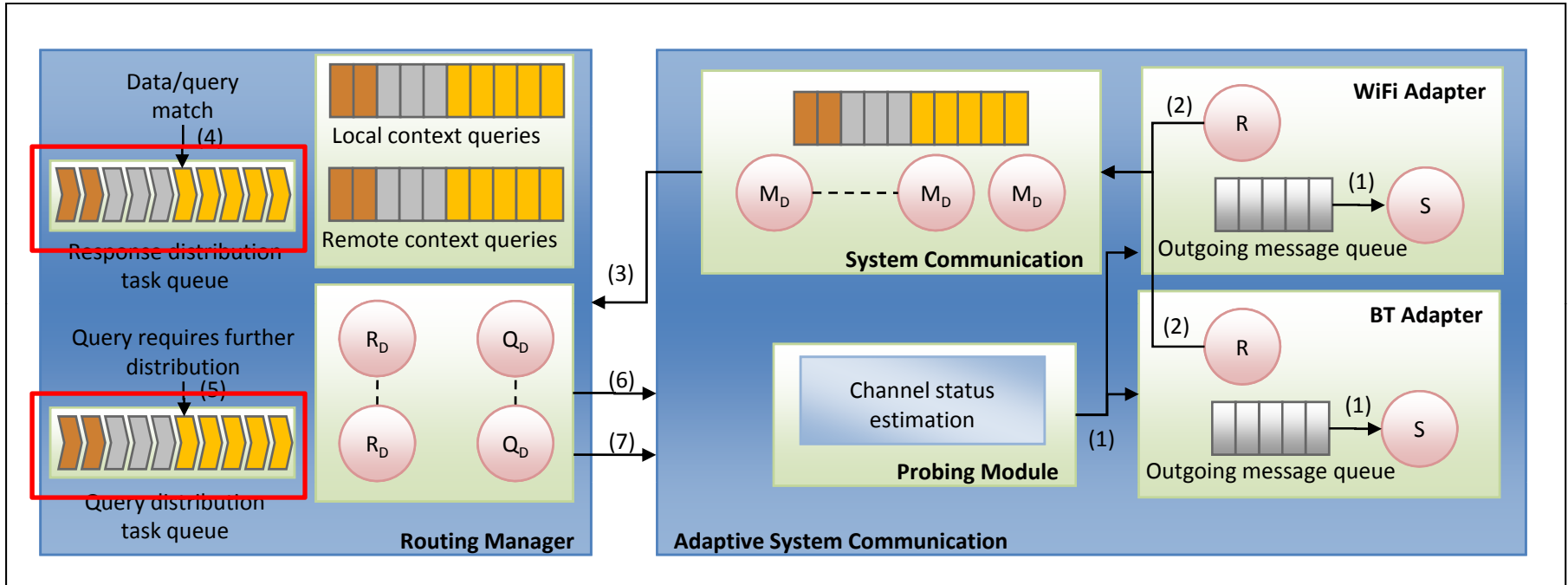
Query Parameters	Depends on	Details
Query Total Lifetime (QTL)	CDDLDA data retrieval time	Equal to the CDDLDA data retrieval time
Query Priority (QP)	CDDLDA priority	Equal to the CDDLDA priority
Horizontal Time To Live (HTTL)	CDDLDA freshness	HTTL = 0, 1, and 2 for respectively gold, silver, and bronze users
Routing Delay (RD)	CDDLDA freshness and data retrieval time	RD is calculated considering the current level and HTTL

Dispatch received queries/data towards the routing manager

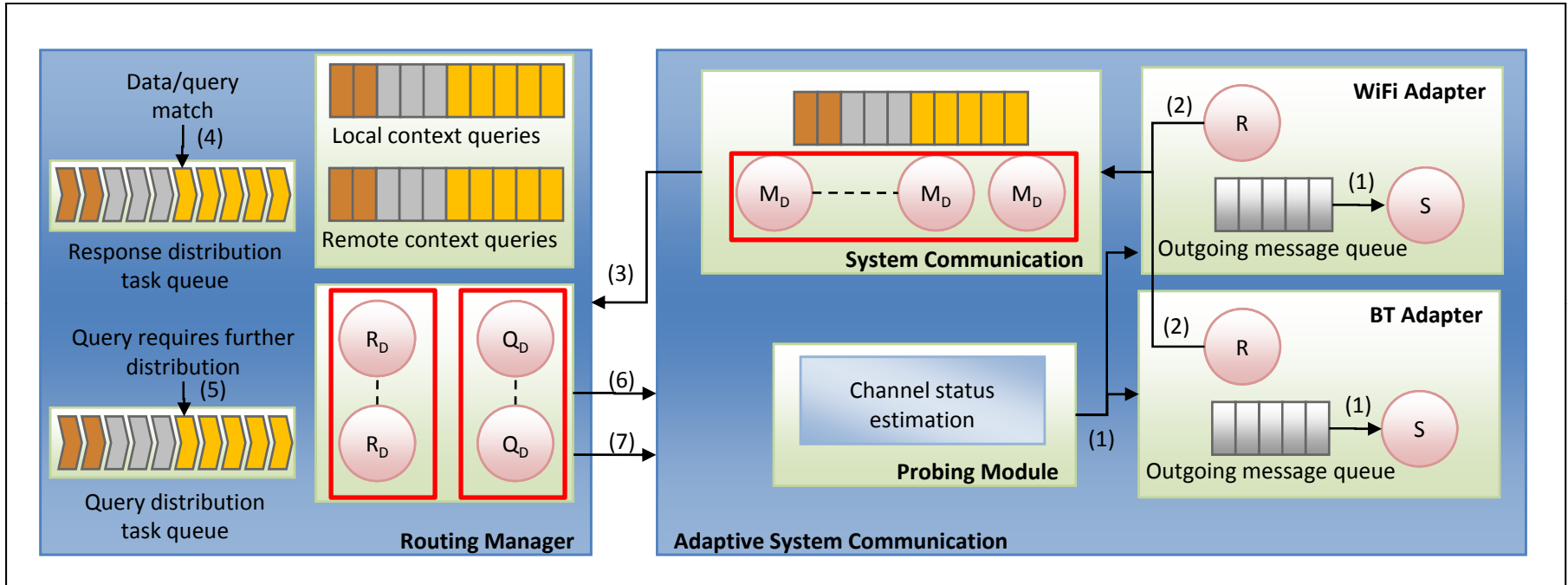


Read querydistribution task queue and send query

- Asynchronous architecture: operations scheduled by means of temporary descriptors
- Minimum intrusion principle → **Limited queue length, Limited processing**

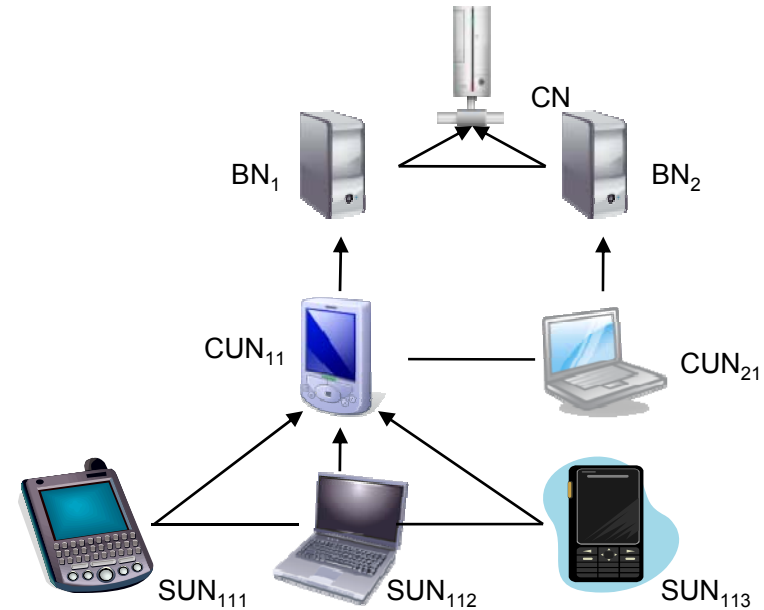


- **Query distribution task** queue ordered according to QP
- **Data distribution task** queue ordered according to the same priority of the matching queries that triggered data distribution
- Queue are partitioned: **50% of queue to priority 0 (gold users), 30% to priority 1 (silver users), and 20% to priority 2 (bronze users)**

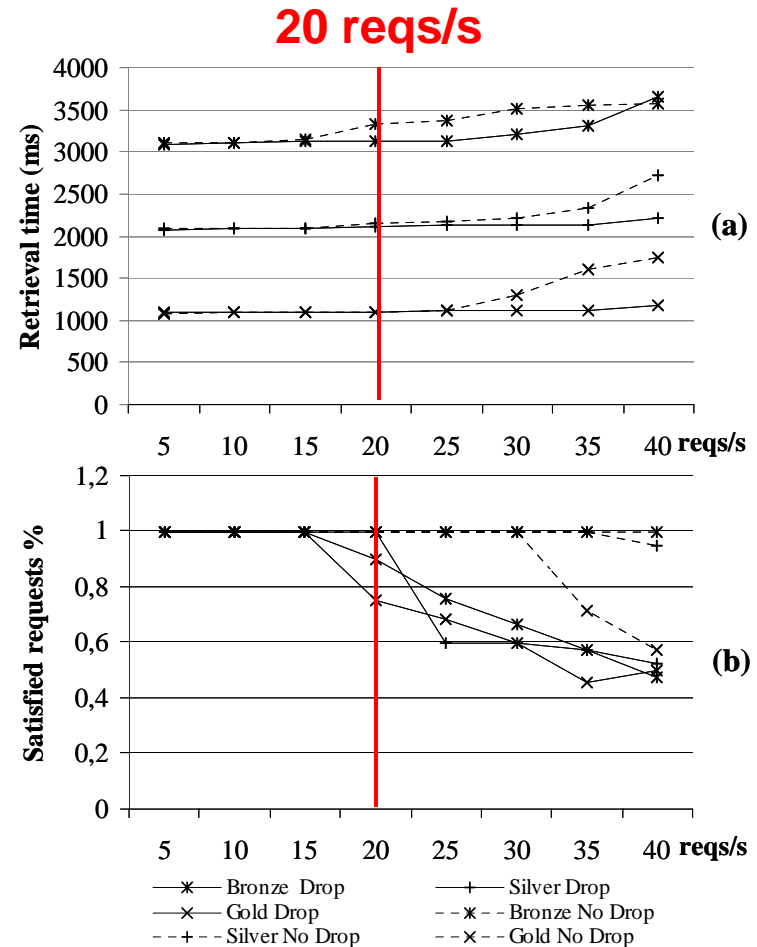


- For each thread pool, we impose number of threads and maximum execution rate for single thread (number of tasks served for second)
- **Proactive approach** to keep queue limited: monitors the time elapsed between the expected (scheduled) task execution time and the effective execution time, and proactively discharge tasks

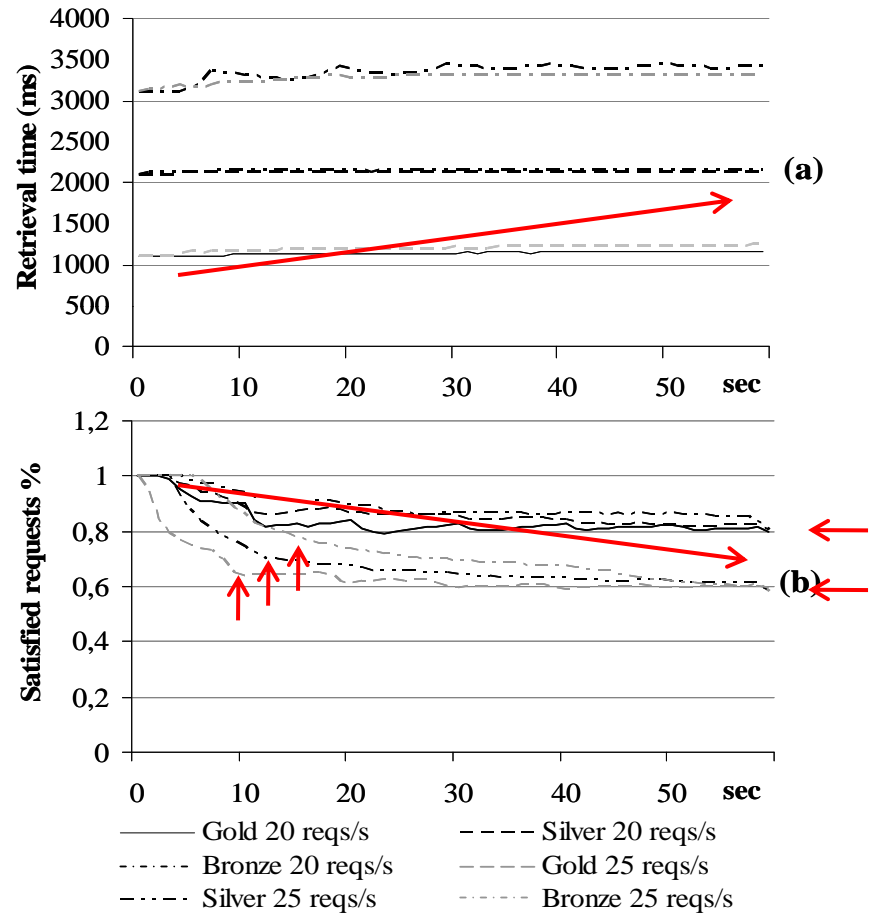
- Implementation insights
  - SALES has been realized on J2SE 1.6
- Experimental testbed
  - CN and BNs execute on 2 CPUs 1,80GHz, 2048MB RAM, Linux Ubuntu
  - Wireless infrastructure composed by Wi-Fi Cisco Aironet 1100 AP
  - Test stations with IEEE 802.11g D-Link WDA-2320 and Linux Ubuntu
- SALES configurations
  - Each SUN belongs to a different user class
  - CUNs/SUNs have 3  $M_D$ , 1  $R_D$ , and 1  $Q_D$  threads; processing rate 20 reqs/s
  - BNs and CN have 10  $M_D$ , 3  $R_D$ , and 3  $Q_D$  threads; processing rate 60 reqs/s
  - All queues are limited to 50 elements for CUNs/SUNs, and to 200 elements for CN/BNs



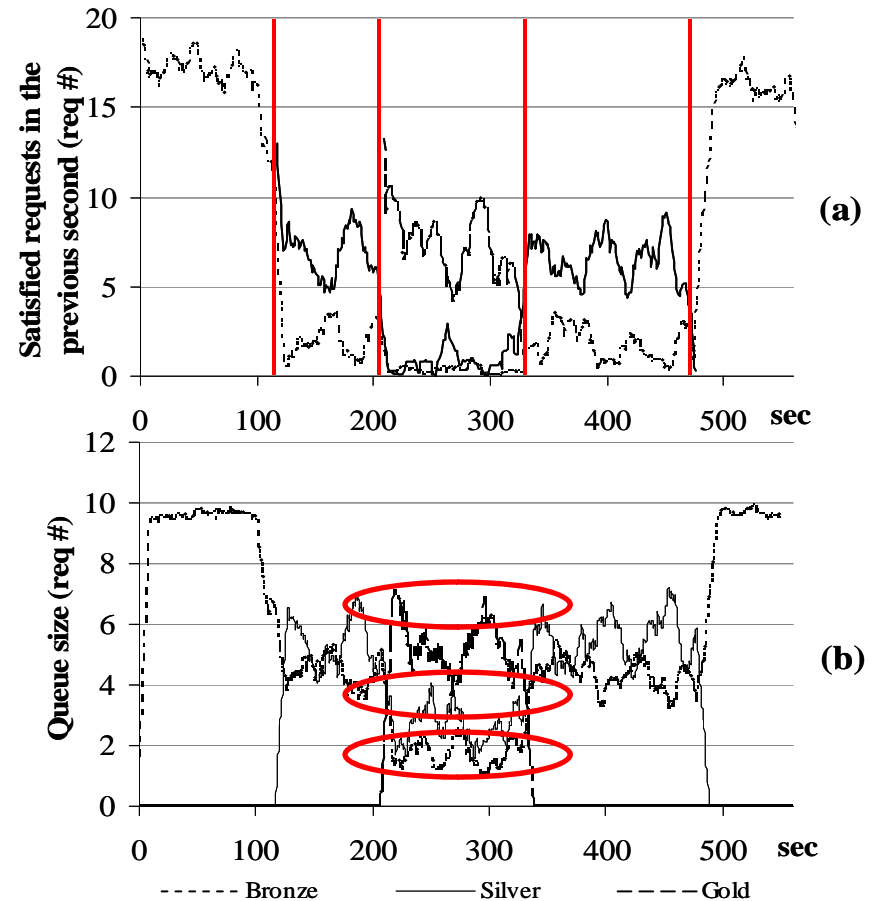
- Query **dropping disabled** (dashed lines) → fewer failures, but higher average retrieval times
- Query **dropping enabled** (continuous lines) → close to the local processing capacity of 20 reqs/s, we experience request failures, but we guarantee limited average response times



- Query dropping enabled
- Worse values during time due to system saturation
- Different clients converge to the same drop rates → SALES exploits all available resources when traffic belongs to only one user class
- Clients belonging to higher classes can experience more failures due to the tight time constraints



- 3 SUNs with different class (gold, silver, bronze); SUNs emit 20 reqs/s
- Each time a SUN of a higher class connects, low-priority SUNs experience higher message drops
- Each time a SUN of a higher class disconnects, low-priority SUNs experience lower message drops
- No one reaches its maximum queue limit → the queue dropping policy does not affect the reported results





## Conclusions

- **CDDLAs** to drive CDDI management and obtain scalability
- **Minimal intrusion** to control CDDI overhead (CPU, memory, ...)
- **Differentiation** to grant QoC (especially under congestion 😊)

## Ongoing work

- **Extensive simulations** to validate our approach in wide deployments composed of several nodes (SALES on ns-2 😊)
- **Extension of the CDDLAs** to introduce data-related quality
- **Different dissemination algorithms**, e.g., flooding- or gossip-based, according to data scope and environmental conditions

Prototype code and information:

**<http://lia.deis.unibo.it/Research/SALES>**

Contacts: Mario Fanelli (mario.fanelli@unibo.it)

Thanks for your attention!